Lecture 4.3: Closures and Equivalence Relations

CS 250, Discrete Structures, Fall 2011

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*Adopted from previous lectures by Cinda Heeren

Course Admin

- Mid-Term 2 Exam Graded
 - Solution has been posted
 - Any questions, please contact me directly
 - Will distribute today
- HW3 Graded
 - Solution has been posted
 - Please contact the TA for questions
 - Will distribute today

Final Exam

- Thursday, December 8, 10:45am-1:15pm, lecture room
 - Heads up!
 - Please mark the date/time/place
- Our last lecture will be on December 6
 We plan to do a final exam review then

HW4

- Expect HW4 to be posted later this week
- Covers the chapter on Relations
- Will be due in 10 days from the day of posting

Outline

Closures

Equivalence Relations

N-ary Relations

- So far, we were talking about binary relations
 defined on two sets.
- Can be generalized to N sets
- Ex: R = {(a, b, c): a < b < c}, defined on set of integers – a 3-ary relation
- Applications in databases

Closure

- Consider relation R={(1,2),(2,2),(3,3)} on the set A = {1,2,3,4}.
- Is R reflexive?



What can we add to R to make it reflexive? (1,1), (4,4)

R' = R U {(1,1),(4,4)} is called the *reflexive closure* of R.

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Lecture 4.3 -- Closures and Equivalence Relations

Closure

Definition:

The closure of relation R on set A with respect to property P is the relation R' with

- $\mathbf{R} \subseteq \mathsf{R}'$
- 2. R' has property P
- 3. \forall S with R \subseteq S and S has property P, R' \subseteq S.

Reflexive Closure

 Let r(R) denote the reflexive closure of relation R.

Then $r(R) = R \cup \{ (a,a): \forall a \in A \}$

- Fine, but does that satisfy the definition?
 - 1. $\mathsf{R} \subseteq \mathsf{r}(\mathsf{R})$ We added edges!
 - 2. r(R) is reflexive By defn
 - 3. Need to show that for any S with particular properties, $r(R) \subseteq S$.

Let S be such that $\mathsf{R} \subseteq \mathsf{S}$ and S is reflexive. Then

$$\{ (a,a): \forall a \in A \} \subseteq S \text{ (since S is reflexive) and } R \subseteq S \\ (given). So, r(R) \subseteq S_{AB} - Closures and Equivalence} \\ Relations$$

Symmetric Closure

 Let s(R) denote the symmetric closure of relation R.

Then $s(R) = R \cup \{ (b,a): (a,b) \in R \}$

- Fine, but does that satisfy the definition?
 - 1. $R \subseteq s(R)$ We added edges!
 - 2. s(R) is symmetric By defn
 - 3. Need to show that for any S with particular properties, $s(R) \subseteq S$.

Let S be such that $R \subseteq S$ and S is symmetric. Then

 $\{ (b,a): (a,b) \in R \} \subseteq S \text{ (since S is symmetric) and } R \subseteq S \\ (given). So, s(R) \subseteq S \\ \underset{Relations}{ So, s(R) } S \in S \\ \underset{Relations}{ So, s(R) } \\ \underset{Relations}{ } \\ \underset{Relations}{ So, s(R) } \\ \underset{Relations}{ }$

Transitive Closure

 Let t(R) denote the transitive closure of relation R.

Then t(R) = RU { $(a,c): \exists b (a,b), (b,c) \in R$

Example: A={1,2,3,4}, R={(1,2),(2,3),(3,4)}.
Apply definition to get:
t(R) = {(1,2),(2,3),(3,4), (1,3), (2,4), (1,3),

Which of the following is true:

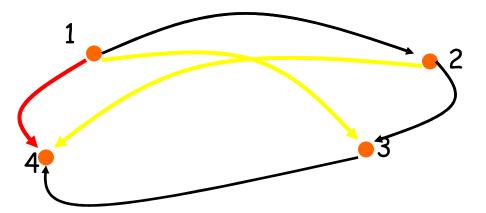
- a) This set is transitive, but we added too much.
- b) This set is the transitive closure of R.
- c) This set is not transitive, one pair is missing.
- d) This set is not transitive, more than 1 pair is missing.

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Relations

Transitive Closure

- So how DO we find the transitive closure?
- Example: A={1,2,3,4}, R={(1,2),(2,3),(3,4)}.



Define a *path* in a relation R, on A to be a sequence of elements from A: $a_i x_{1,...} x_{i,...} x_{n-1}$, b, with $(a_i, x_1) \in R$, $\forall i (x_i, x_{i+1}) \in R$, $(x_{n-1}, b) \in R$.

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"Path from a to b."

Draw a graph.

Transitive Closure

Formally:

If t(R) is the transitive closure of R, and if R contains a path from a to b, then (a,b) \in t(R)

A technique:

- For a set R consisting of n elements, t(R) can be specified by the matrix: $M_R V M_{R^2} V \dots V M_{R^n}$
- More efficient method: Warshall's algorithm

Transitive Closure -- Example

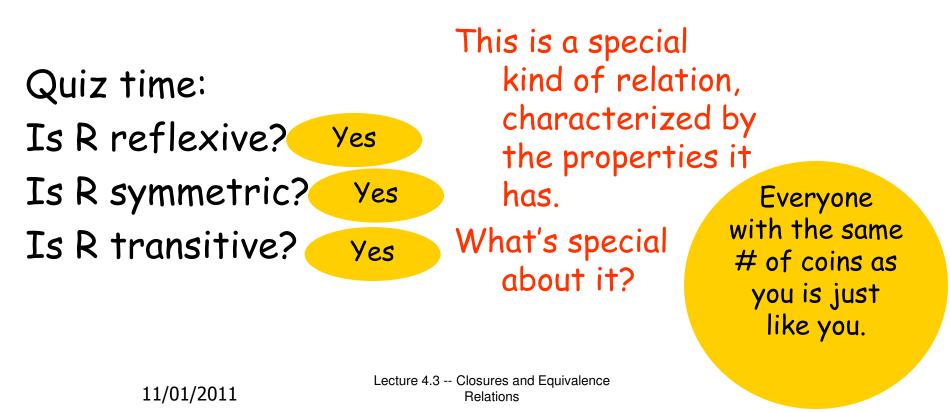
$$M_{R} = \begin{pmatrix} 1 & 1 & 1 \\ 0 & 1 & 1 \\ 1 & 1 & 1 \end{pmatrix}$$

$$\bullet M_R V M_{R^2} V \dots V M_{R^n}?$$

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Example:

Let S = {people in this classroom}, and let R = {(a,b): a has same # of coins in his/her bag as b}



Formally:

Relation R on A is an *equivalence relation* if R is

- Reflexive (∀ a ∈ A, aRa)
- Symmetric (aRb --> bRa)
- Transitive (aRb AND bRc --> aRc)

Recall:

aRb denotes (a,b) ∈ R.

Example:

S = Z (integers), $R = \{(a,b) : a \equiv b \mod 4\}$

Is this relation an equivalence relation on S?

Have to PROVE reflexive, symmetric, transitive.

Example:

S = Z (integers), R = {(a,b) : $a \equiv b \mod 4$ }

Is this relation an equivalence relation on S?

Start by thinking of R a different way: aRb iff there is an int k so that a = 4k + b. Your quest becomes one of finding ks.

Let a be any integer. aRa since a = 4.0 + a. Consider aRb. Then a = 4k + b. But b = -4k + a. Consider aRb and bRc. Then a = 4k + b and b = 4j + c. So, a = 4k + 4j + c = 4(k+j) + c.

Example:

- S = Z (integers), R = {(a,b) : a = b or a =- b}.
 Is this relation an equivalence relation on S?
- Have to prove reflexive, symmetric, transitive.

Example:

- S = R (real numbers), R = {(a,b) : a b is an integer}. Is this relation an equivalence relation on S?
- Have to prove reflexive, symmetric, transitive.

Example:

- S = R (real numbers), R = {(a,b) : a b is an integer}. Is this relation an equivalence relation on S?
- Have to prove reflexive, symmetric, transitive.

- Example:
- S = N (natural numbers), R = {(a,b) : a | b}. Is this relation an equivalence relation on S?
- Have to prove reflexive, symmetric, transitive.

Equivalence Classes

Example: Back to coins in bags.

Definition: Let R be an equivalence relation on S. The equivalence class of $a \in S$, $[a]_R$, is $[a]_R = \{b: aRb\}$

a is just a name for the equiv class. Any member of the class is a representative.

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Equivalence Classes

Definition: Let R be an equivalence relation on S. The equivalence class of $a \in S$, $[a]_R$, is

 $[a]_R = \{b: aRb\}$

Notice this is just a subset of S.

What does the set of equivalence classes on S look like?

To answer, think about the relation from before:

- S = {people in this room}
- R = {(a,b) : a has the same # of coins in his/her bag as b}
- In how many different equivalence classes can each person fall?

Equivalence Relations

Similarly, consider the $a \equiv b \mod 4$ relation Lecture 4.3 -- Closures and

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Rosen 9.4 and 9.5

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